



RAMA UNIVERSITY

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FACULTY OF ENGINEERING & TECHNOLOGY

BCS-501 Operating System

Lecturer-23

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Assistant Professor

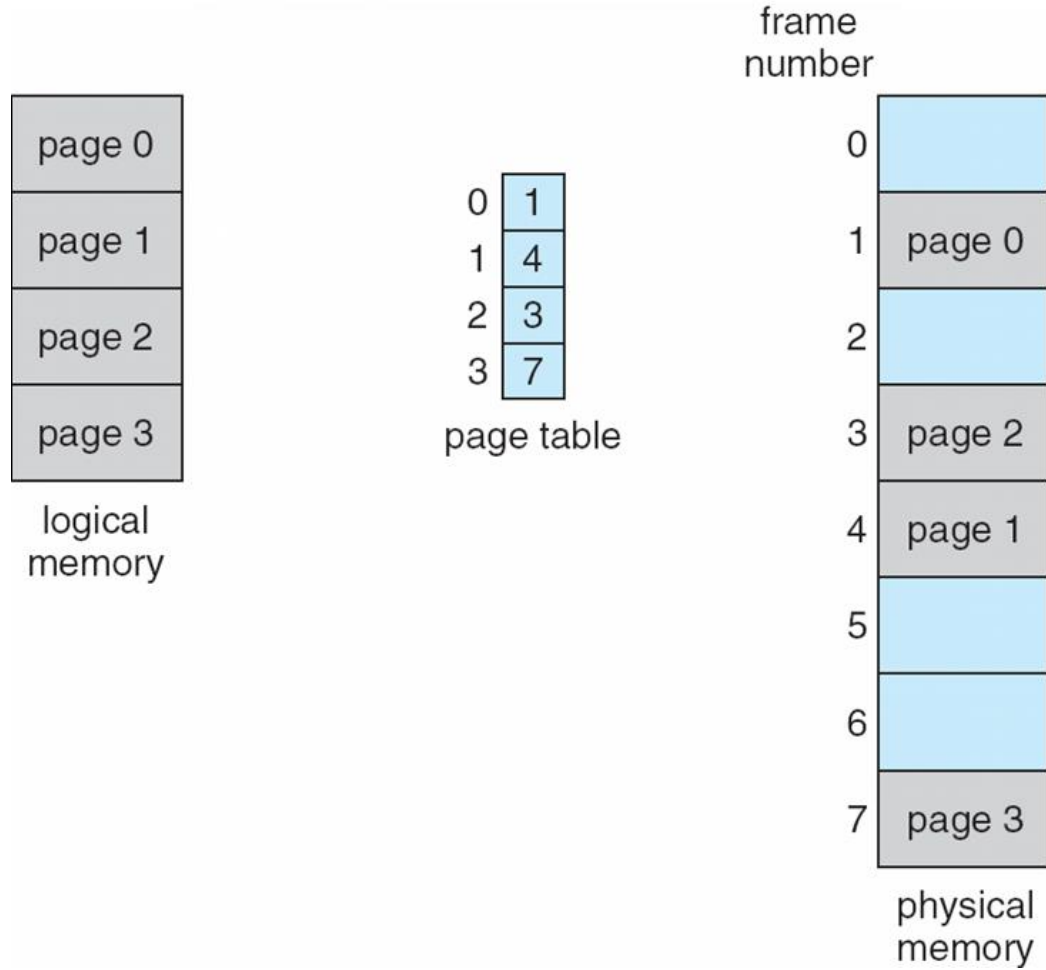
Computer Science & Engineering

Memory

- **Paging Model of Logical and Physical Memory**
- **Paging Example**
- **Frames**



Paging Model of Logical and Physical Memory



Paging Example

0	a
1	b
2	c
3	d
4	e
5	f
6	g
7	h
8	i
9	j
10	k
11	l
12	m
13	n
14	o
15	p

logical memory

0	5
1	6
2	1
3	2

page table

0	
4	i j k l
8	m n o p
12	
16	
20	a b c d
24	e f g h
28	

physical memory

Paging (Cont.)

- Calculating internal fragmentation—

Page size = 2,048 bytes

Process size = 72,766 bytes

35 pages + 1,086 bytes

Internal fragmentation of 2,048 - 1,086 = 962 bytes

Worst case fragmentation = 1 frame – 1 byte

On average fragmentation = 1 / 2 frame size

- So small frame sizes desirable?

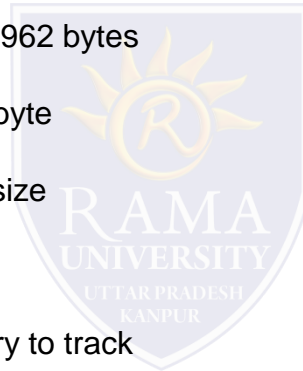
- But each page table entry takes memory to track

- Page sizes growing over time

 - Solaris supports two page sizes – 8 KB and 4 MB

- Process view and physical memory now very different

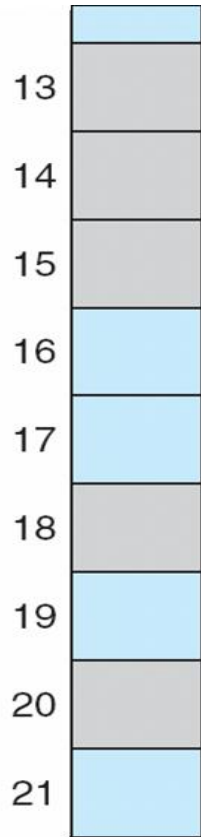
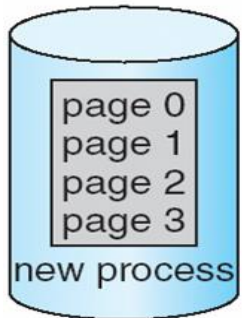
- By implementation process can only access its own memory



Free Frames

free-frame list

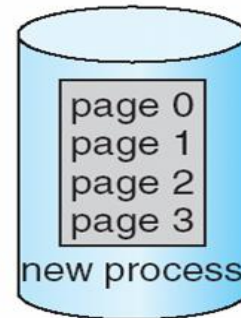
14
13
18
20
15



(a)

free-frame list

15



0	14
1	13
2	18
3	20

new-process page table



(b)

Paging is implemented in

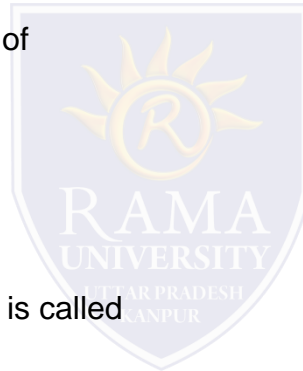
- A. Operating System
- B. Hardware
- C. Software
- D. All of them

Page-Table length register (PTLR) indicates size of

- A. Page Table
- B. Paging File
- C. Main Memory
- D. Virtual Memory

Bring a page into memory only when it is needed is called

- A. Demand Memory
- B. Demand Paging
- C. Page Fault
- D. Page Segmentation



Bring a page into memory only when it is needed is called

- A. Demand Memory
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Dived logical memory into blocks with the same size as frames are called

- A. Pages
- B. Frames
- C. Page Table
- D. Segmentation

